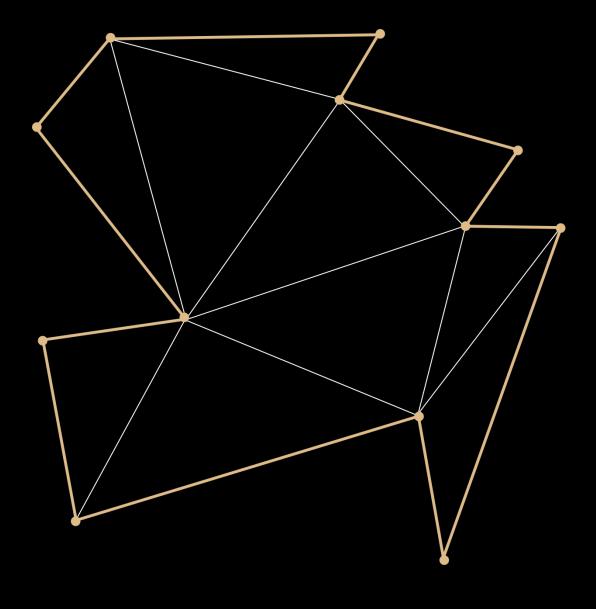
Rectangular Spiral Galaxies are Still Hard



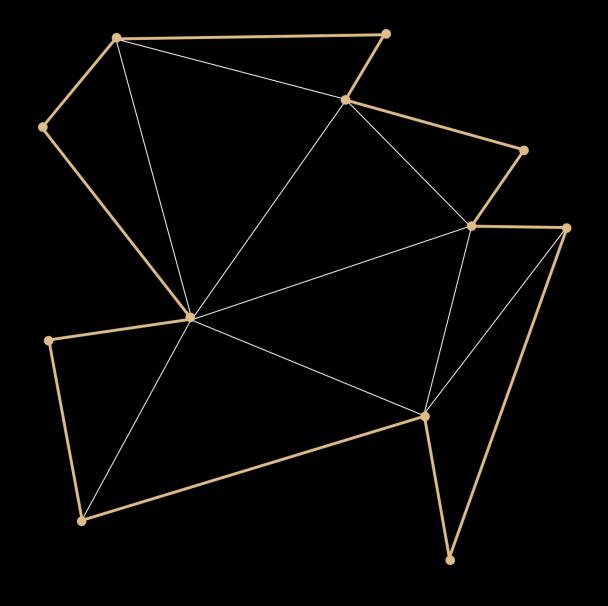




★Triangulation



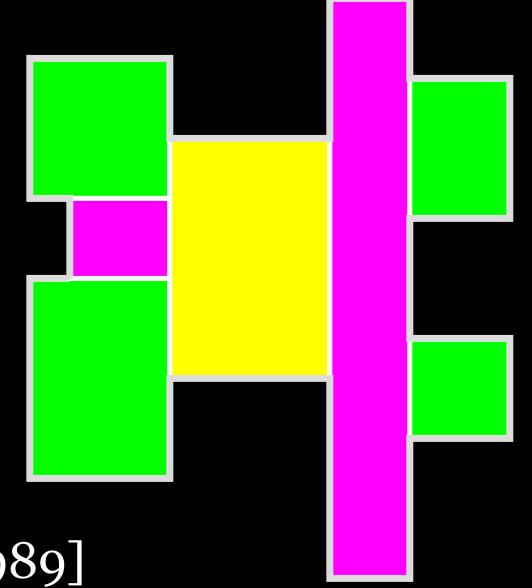
- **★**Triangulation
- ★ Partition rectangular polygons into rectangles







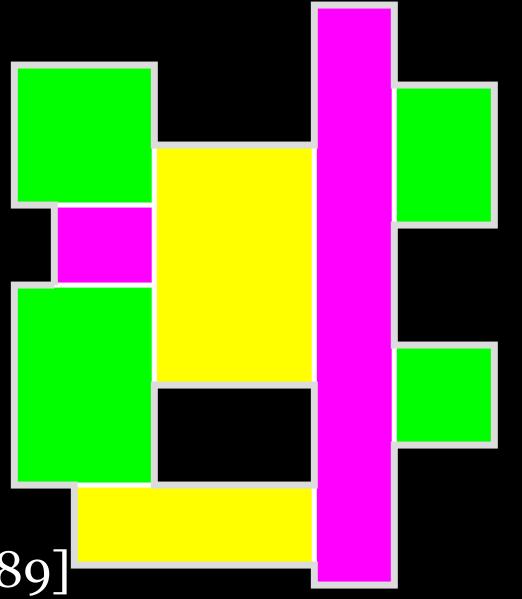
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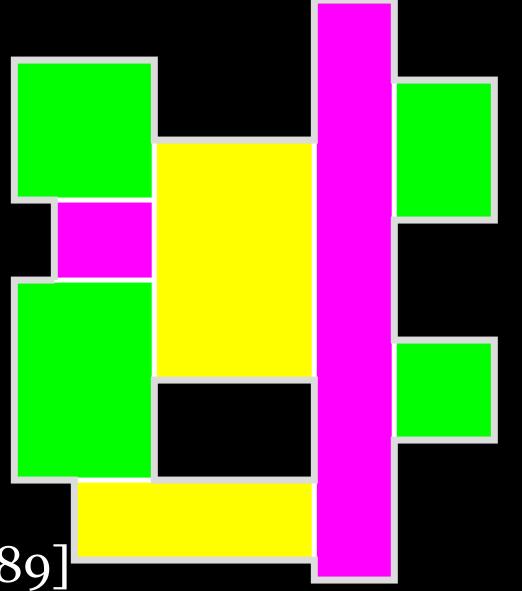
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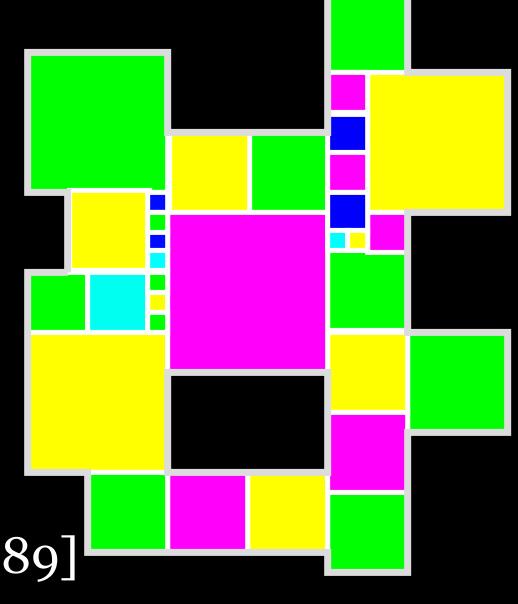
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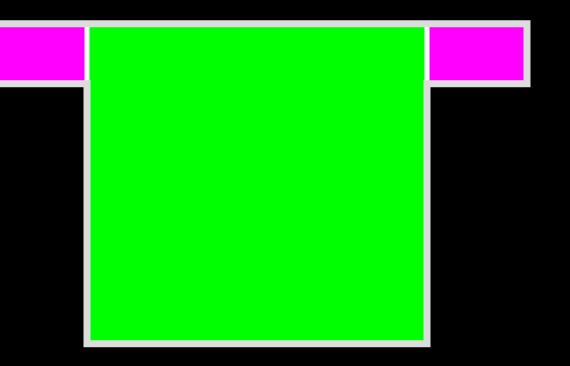




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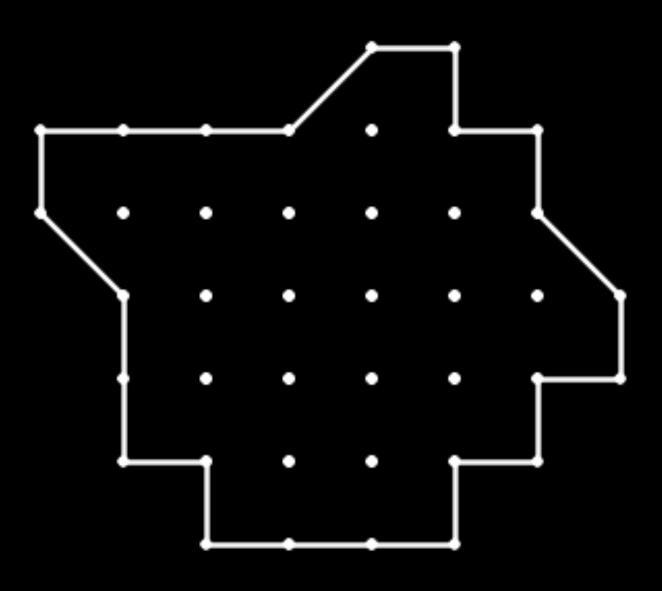


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- **★** Partition rectangular polygons into polygons with at most 8 vertices: [3n+4/16] polygons with O(n) algorithm [Győri, Mezei, 2016]





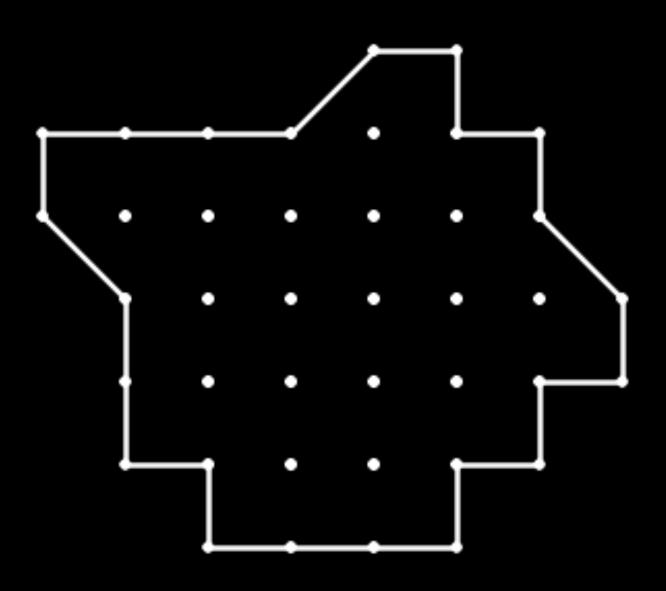
2







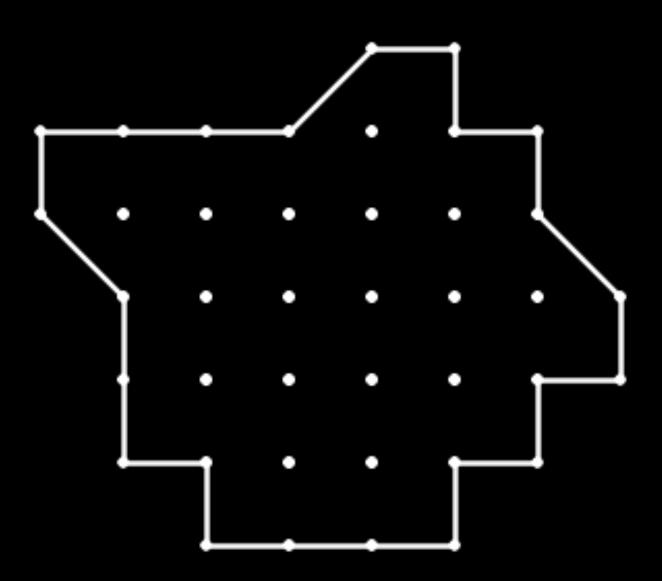
★ Partition into "similar" polygons:







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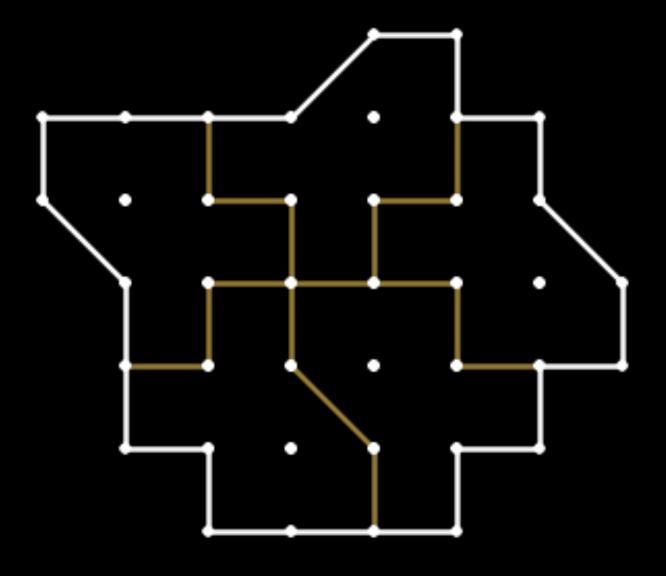






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Several versions, most common: homothetic—same size, can mirror

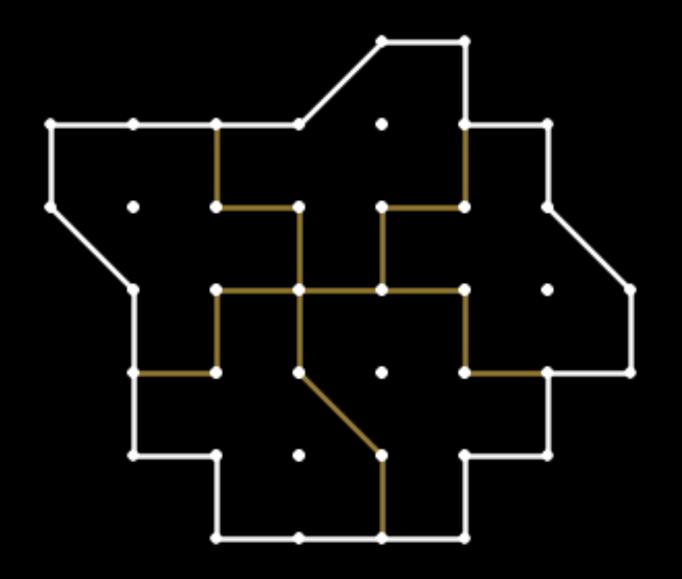






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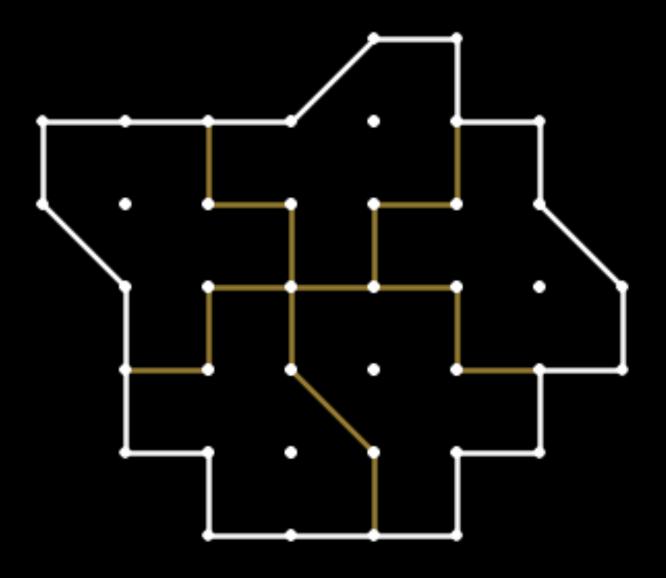






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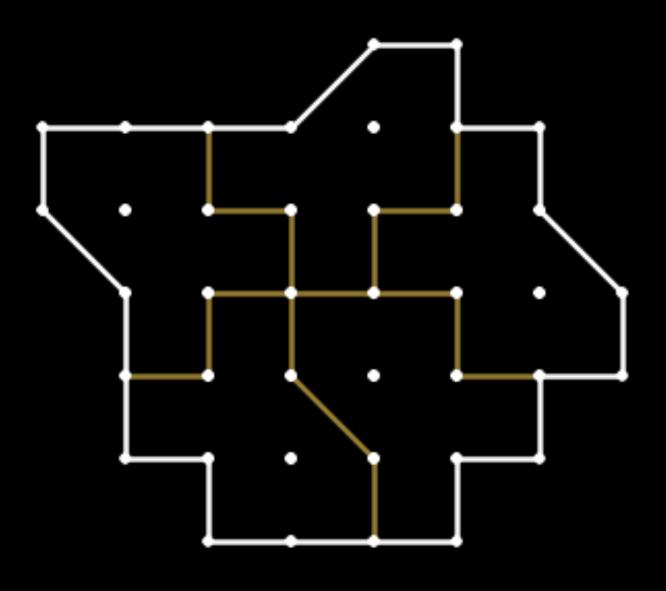






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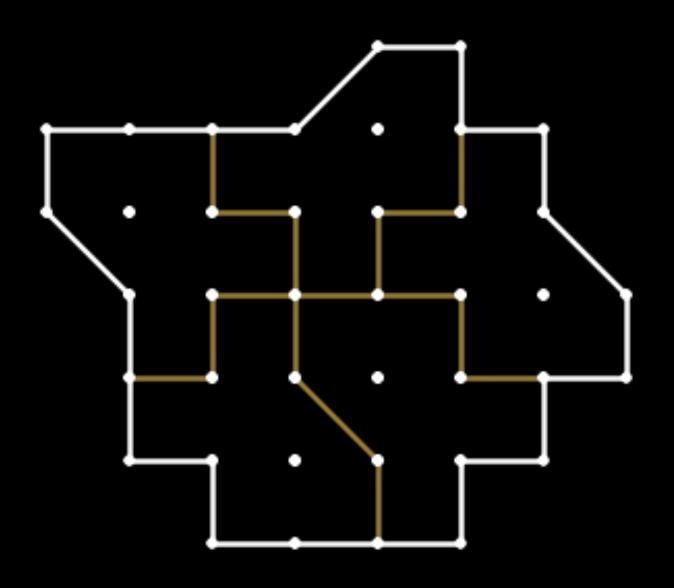






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- ★ Rectangular polygons → let's look at polyominoes

Several versions, most common: homothetic—same size, can mirror













★ Pencil-and-paper puzzle from Nikoli ("Tentai Show")

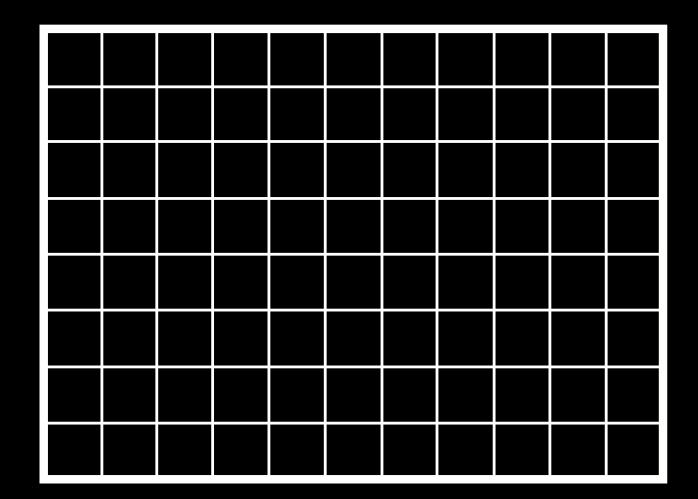






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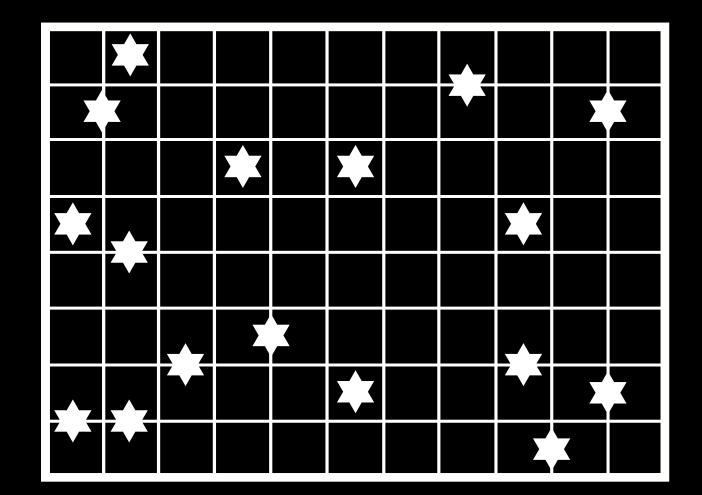






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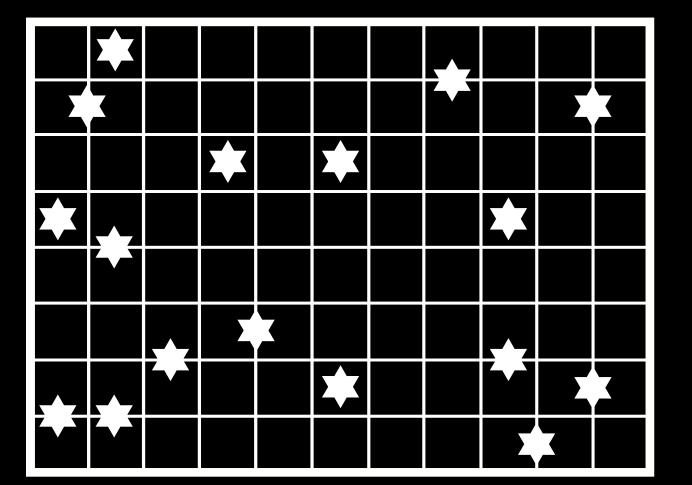






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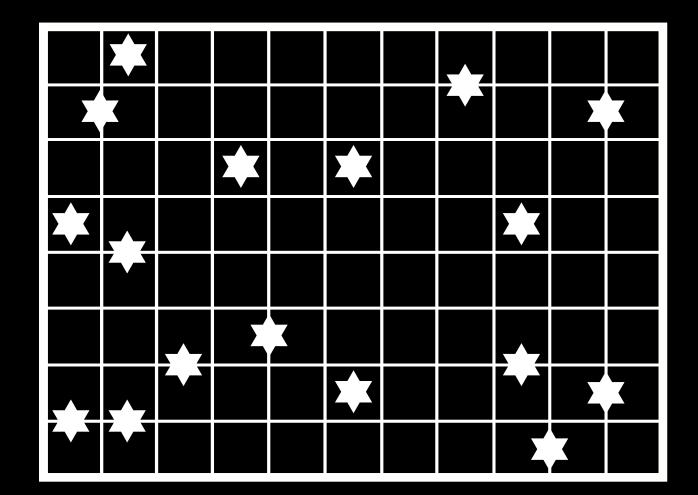






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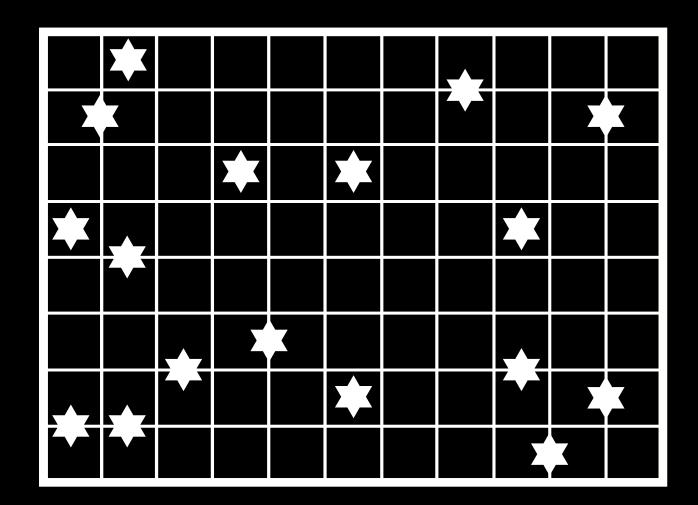






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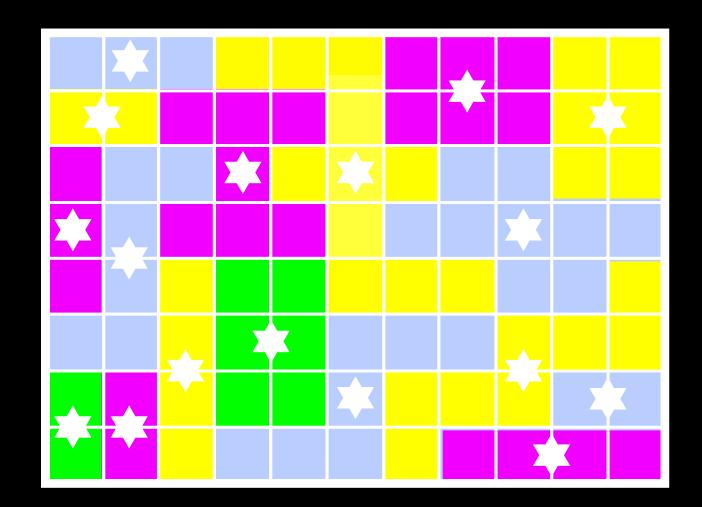






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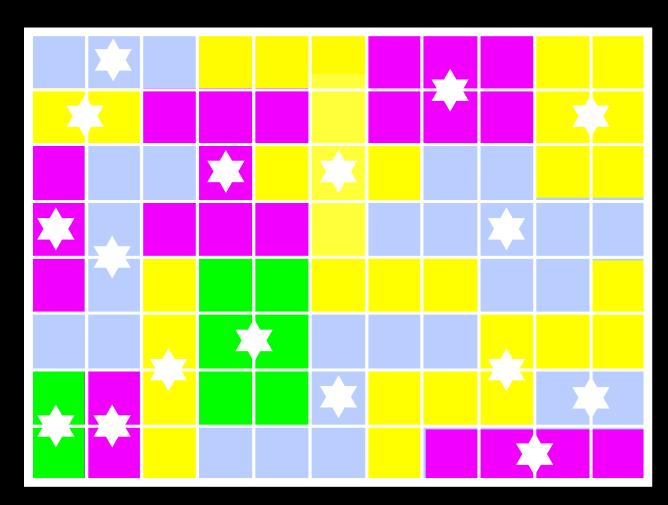






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Also gives us a tiling using the solution galaxies as polyominoes





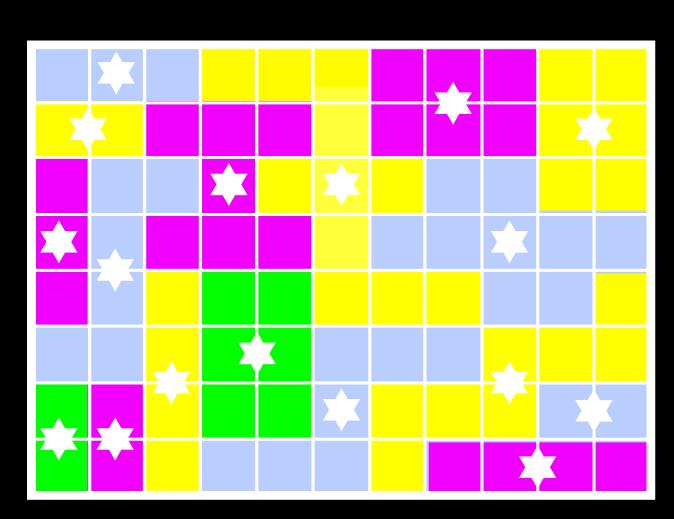
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- ★ Do you want to solve some puzzles?

http://www.nikoli.co.jp/en/puzzles/astronomical_show.html

http://puzzlepicnic.com/genre?id=17

https://www.gmpuzzles.com/blog/spiral-galaxies-rules-info/





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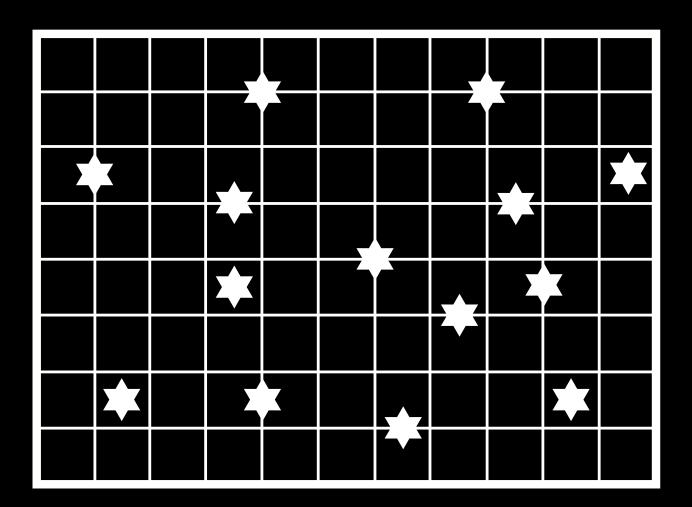






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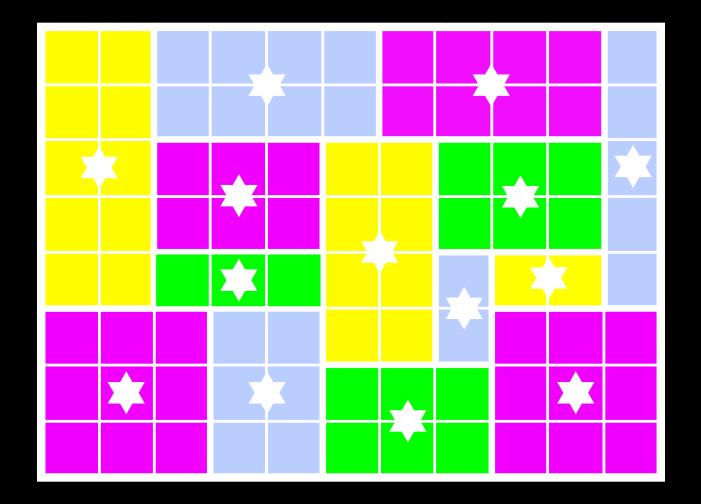






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Generating Spiral Galaxies Puzzles







Generating Spiral Galaxies Puzzles

* What do we want to optimize when we generate a puzzle?







Generating Spiral Galaxies Puzzles

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- ★ Interesting optimization question for a coloured version:

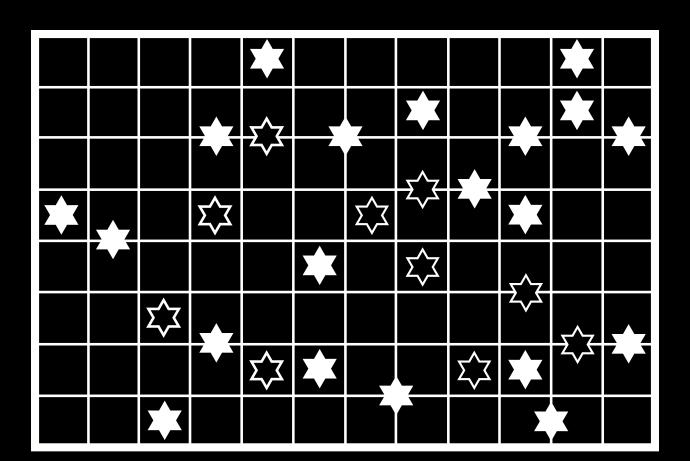






- ★ What do we want to optimize when we generate a puzzle?
- * Interesting optimization question for a coloured version:
 - * Some centers are coloured black



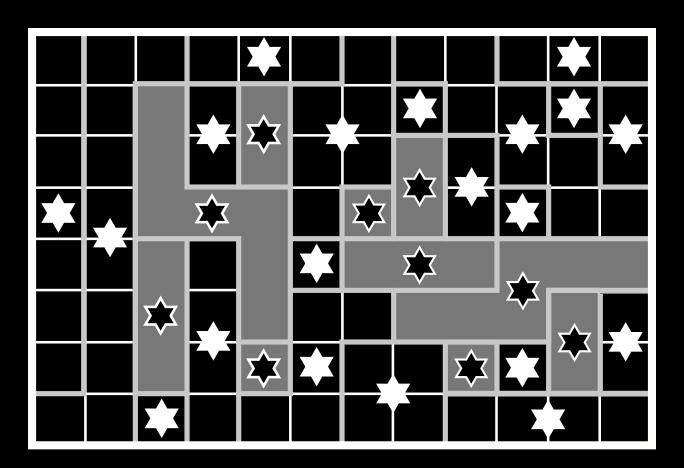






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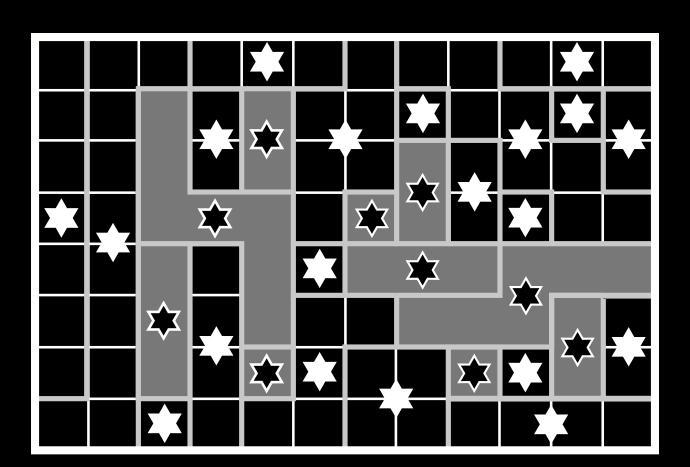






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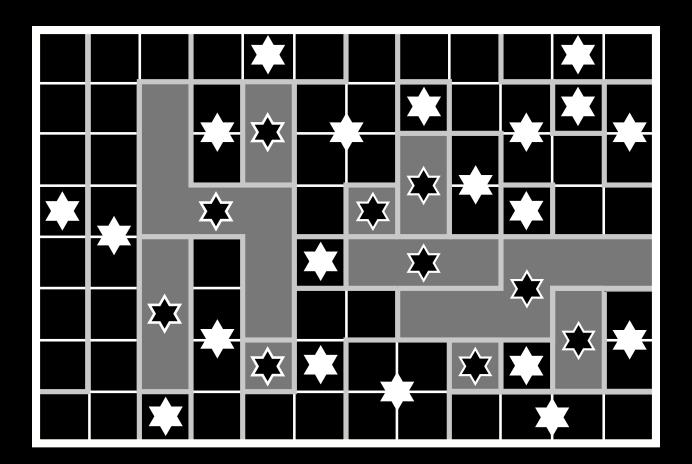






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 - → Minimum number of centers, such that there exist Spiral Galaxies that exactly cover a given shape









Results

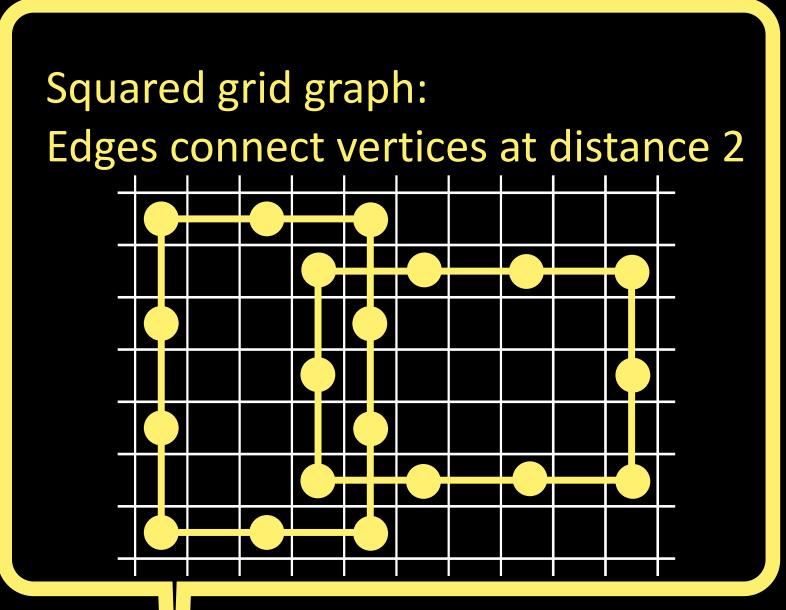
- ★ Determining if a Spiral Galaxies board is solvable with only rectangular galaxies is NP-complete.
- ★ Determining if a Spiral Galaxies board is solvable with only 1x1, 1x3 and 3x1 galaxies is NP-complete and counting the number of solutions is #P-complete and ASP-complete.
- ★ Non-crossing matching in squared grid graphs is NP-complete.
- ★ Generating puzzles: Minimizing the number of centers on a Spiral Galaxies board, such that Spiral Galaxies with these centers exactly cover a given shape S is NP-complete.





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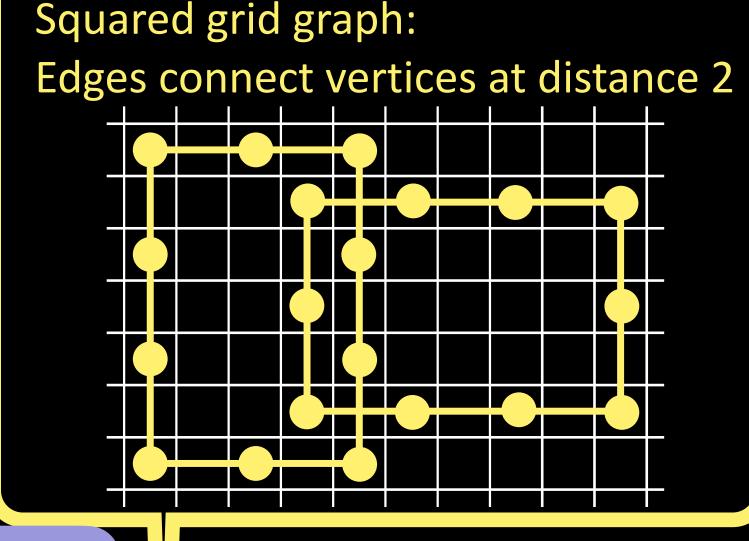
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★ Reduction from PLANAR 1-IN-3 SAT





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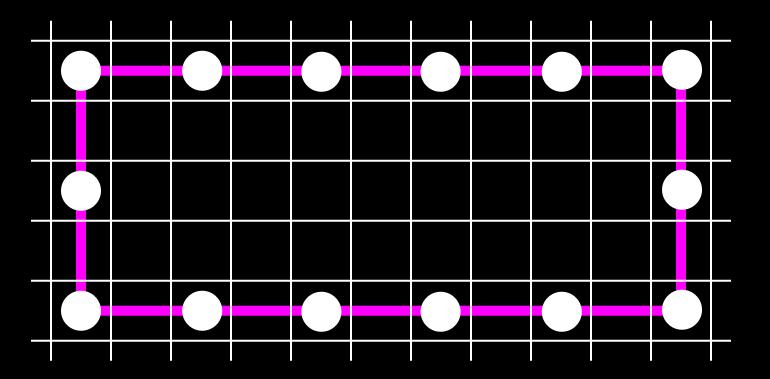


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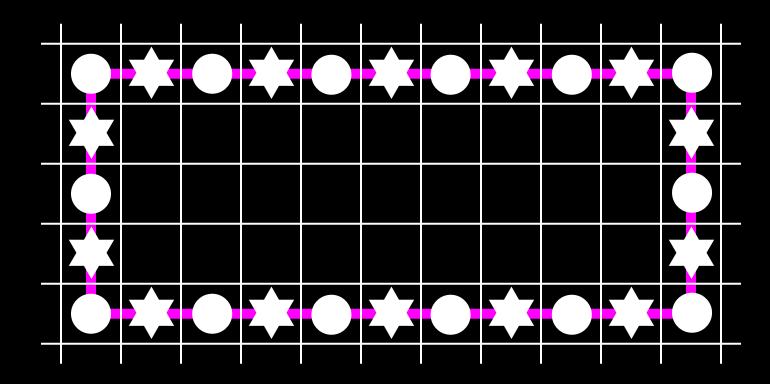
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- ★ Disks with distance 2 can be connected by an edge







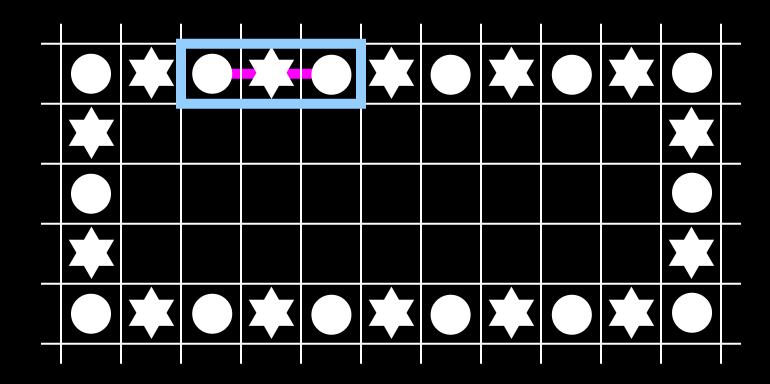
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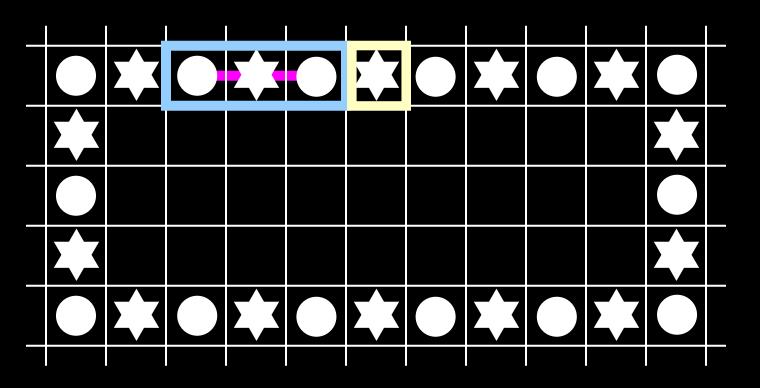
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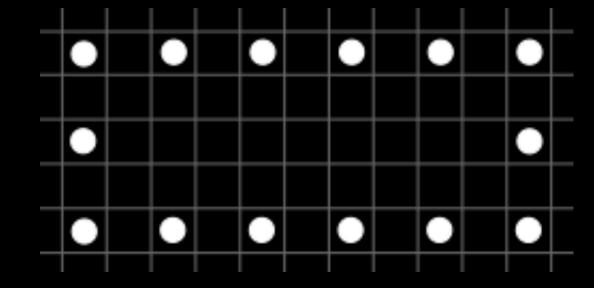
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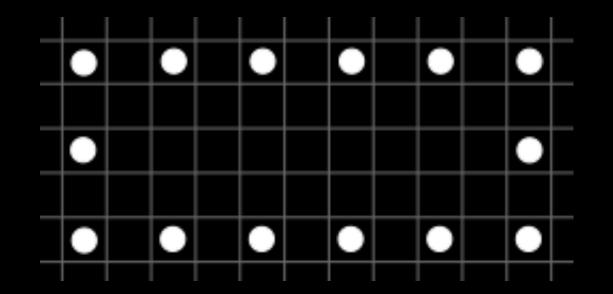




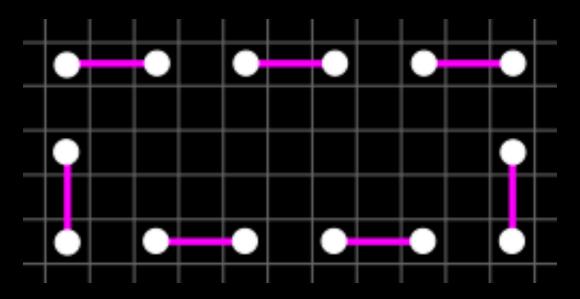
Variable loop

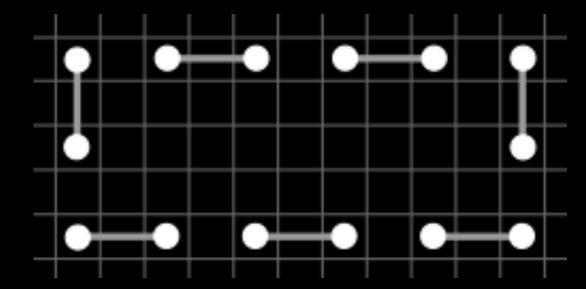


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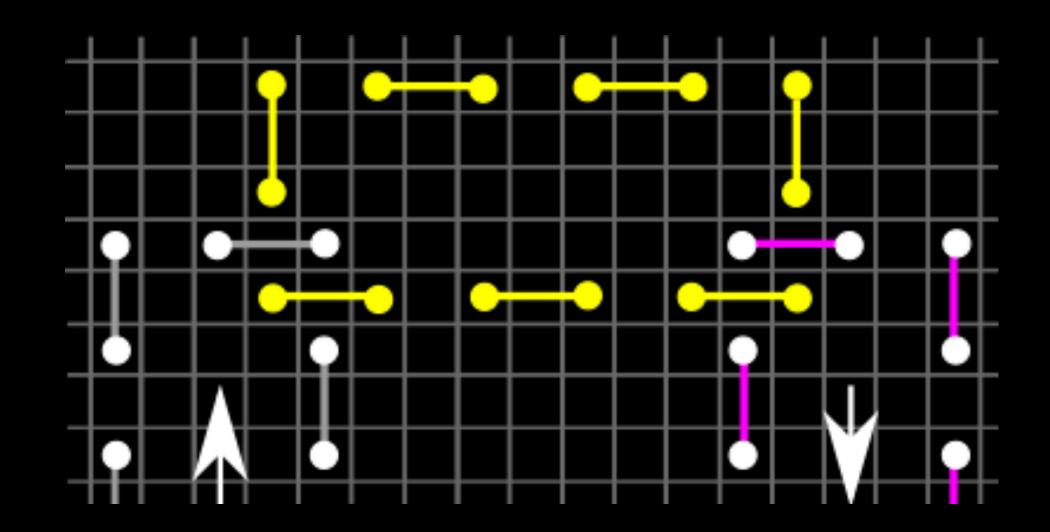
Two possible states—"true" and "false":

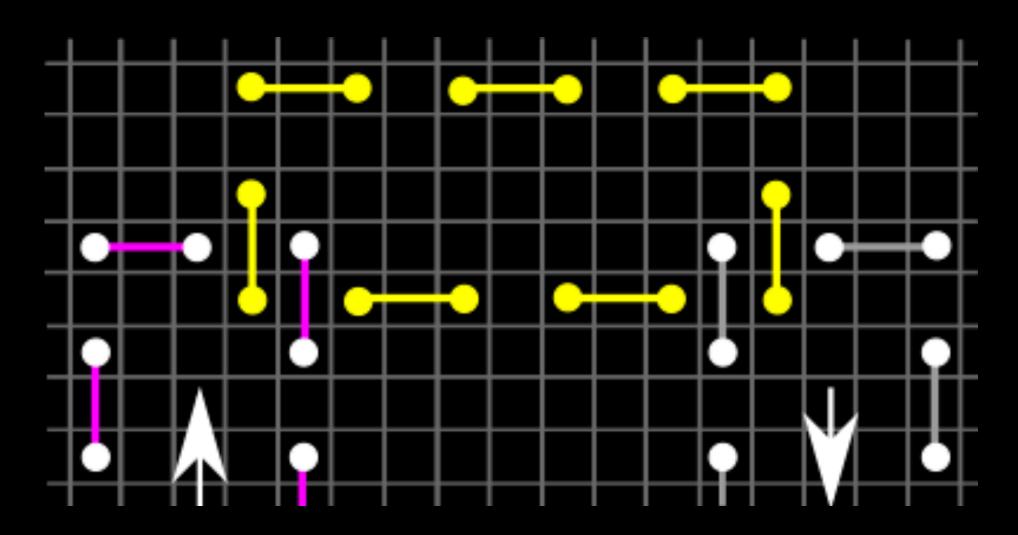






Negation gadget

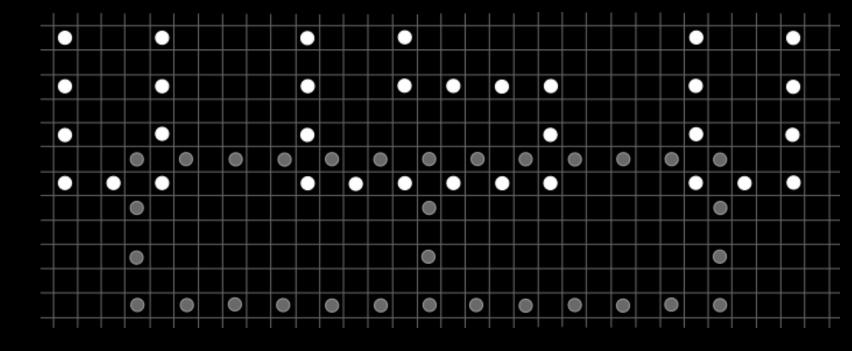






Clause gadget

3 variable loops

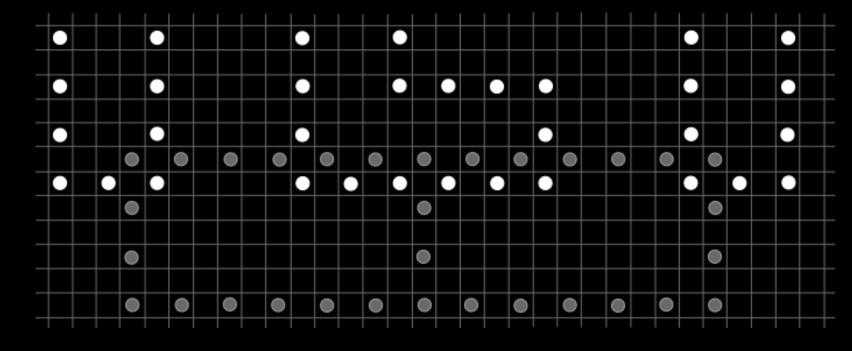


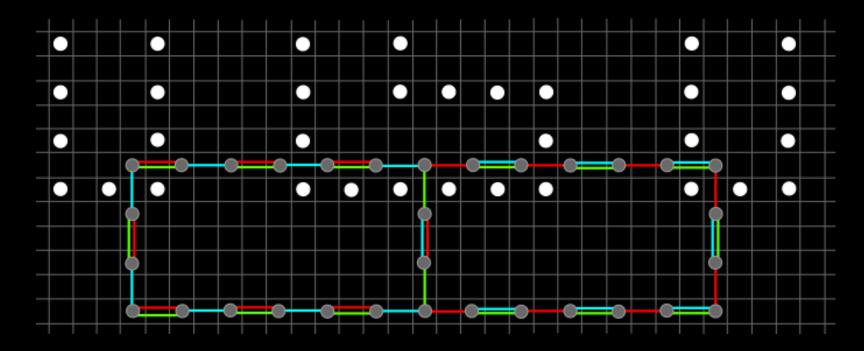




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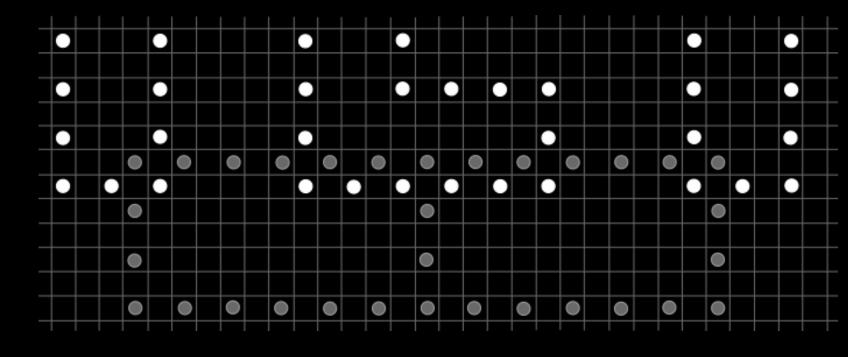


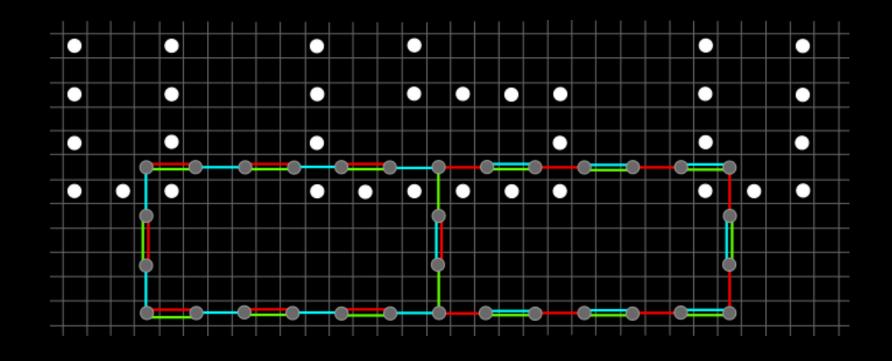


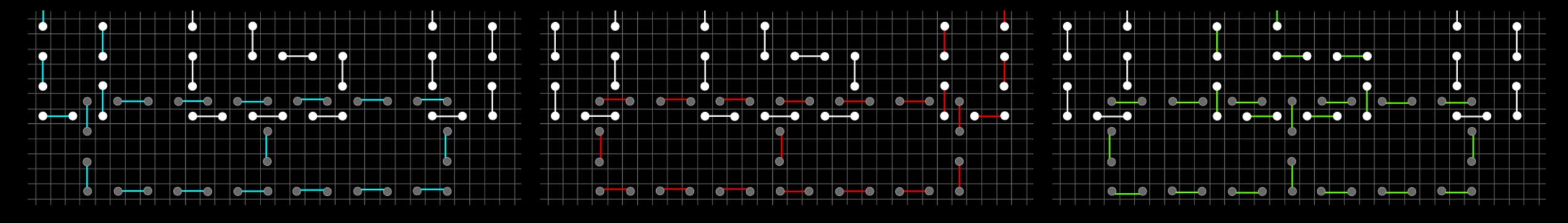


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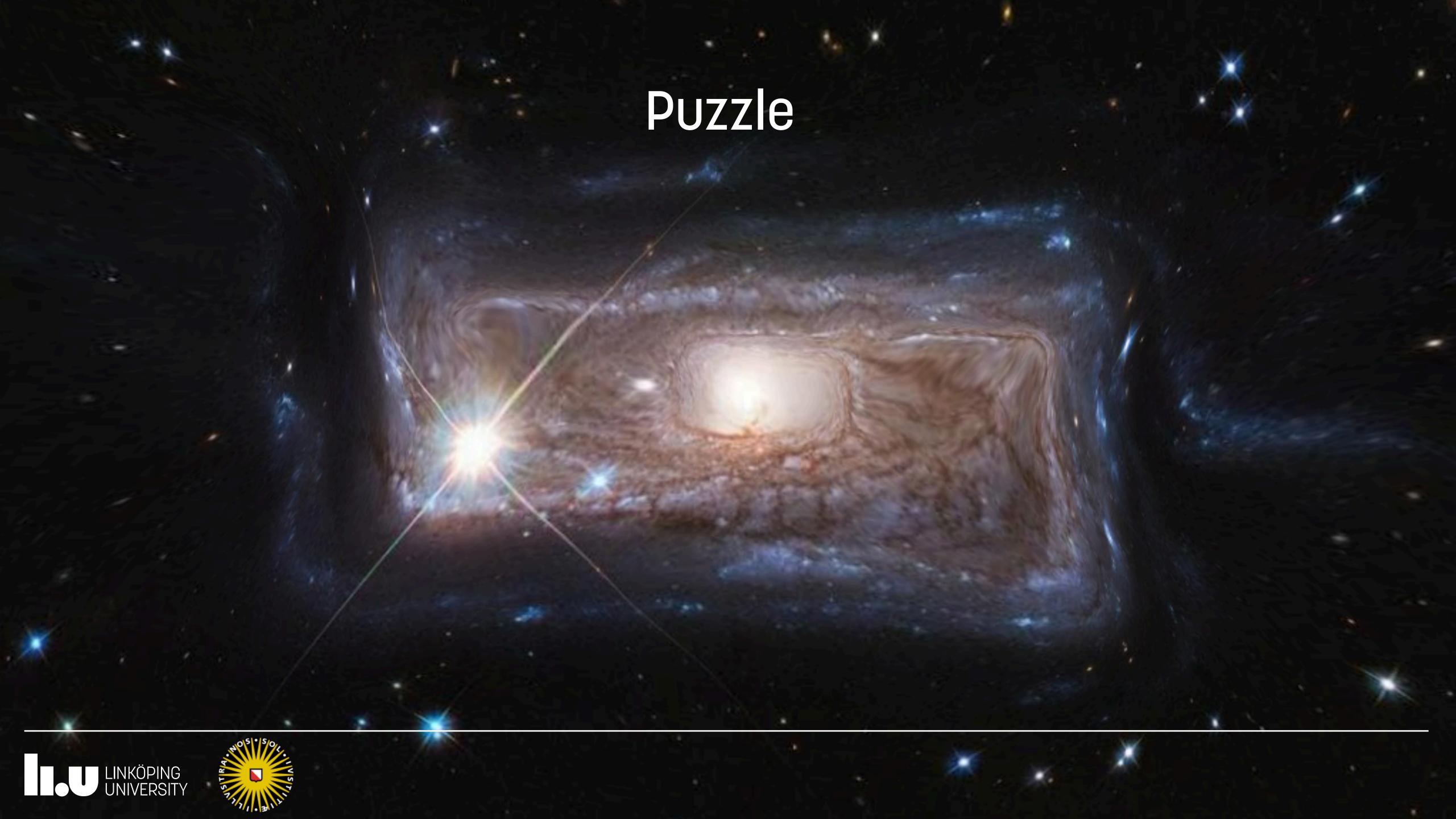




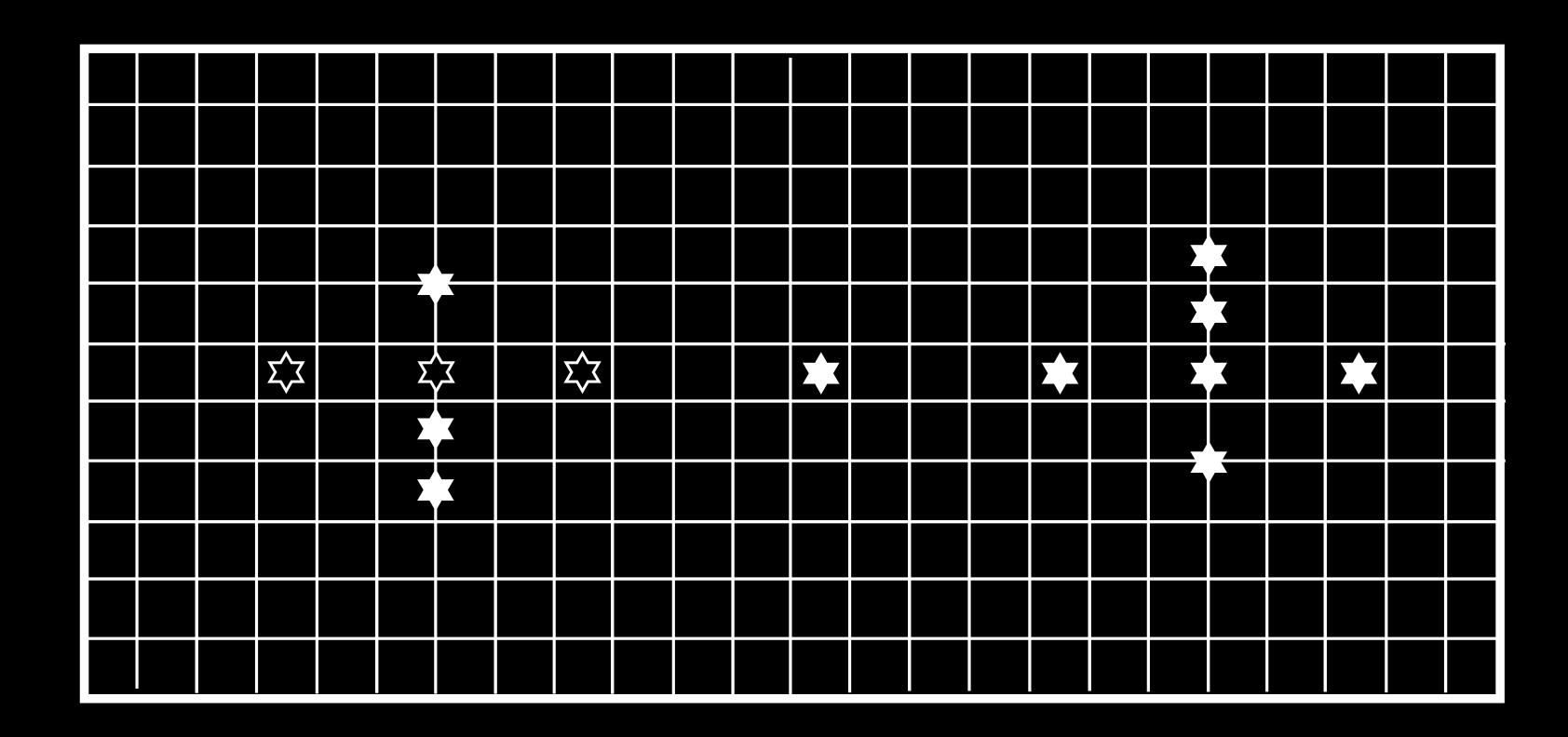








Can be solved for the letters A, B, H, P, R, S, Z (+E for disconnected galaxies)





SCAN ME





